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18DD1905

M. Sc. EXAMINATION, 2020

(Fourth Semester)

(C Scheme)

(Main Only)

MATHEMATICS

MAT612C

ALGEBRAIC CODING THEORY-II

Time: 3 Hours [Maximum Marks: 75]

Before answering the question-paper candidates should ensure that they have been supplied to correct and complete question-paper. No complaint, in this regard, will be entertained after the examination.

Note: Attempt *Five* questions in all, selecting *one* question from each Unit. Q. No. 1 is compulsory. All questions carry equal marks.

(Compulsory Question)

- 1. (a) Write down a parity check matrix for the binary (4, 7) Hamming code.
 - (b) Find a primitive element of the finite field GF(9).
 - (c) Obtain the generator polynomial of the ternary cyclic code $C = \{000, 111, 222\}$.
 - (d) Define Reed-Solomon (RS) Code. Show that RS code is a BCH code.
 - (e) Give an example of a product code.

 $3 \times 5 = 15$

(1-09/28) M-18DD1905

P.T.O.

Unit I

- (a) Prove that every cyclic code of length n over finite field F_q correspondences to an ideal in the ring F_q[x]/<xⁿ 1>. Find the generating polynomial of the ideal in F₂[x]/<x⁴ 1> corresponds to the cyclic code C = {0000, 1010, 0101, 1111}.
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 - (b) Show that every cyclic code has an idempotent generator. Find the generating idempotent for the binary [7, 4]-Hamming code.
- 3. (a) Let C_1 and C_2 be two cyclic codes over \mathbf{F}_q of length n with generating idempotent $e_1(x)$ and $e_2(x)$ respectively. Show that $C_1 + C_2$ and $C_1 \cap C_2$ have generating idemptotents $e_1(x) + e_2(x) e_1(x) e_2(x)$ and $e_1(x) e_2(x)$ respectively.
 - (b) Construct all cyclic codes of length 4 and dimension 2 over finite field \mathbf{F}_5 . How many errors these codes can detect and correct?

Unit II

- 4. (a) Define the cyclotomic coset C_s modulo n over finite field \mathbf{F}_q containing the integer s. If α is a primitive nth root of unity in some extension field of \mathbf{F}_{q} , then show that the minimal polynomial of α^s over \mathbf{F}_q is $\mathbf{\Pi}_{j \in C_s}(x \alpha^j)$.
 - (b) Show that a binary BCH code of length $n = 2^k 1$ and the minimum distance at least d = 2t + 1 can be constructed with check digits at most kt.
- (a) Factorize x¹⁵ -1 over F₂. Hence or otherwie, construct a binary BCH code of length 15 and dimension 8.
 - (b) Let r be an integer with $0 \le r \le m$. Then prove the following:
 - (i) The dimensions of Reed-Muller Code \mathbf{R} (r, m) equals

$$\binom{m}{0} + \binom{m}{1} + \dots + \binom{m}{r}$$

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(ii) The minimum distance of **R** (r, m) equals 2^{m-r} .

Unit III

- 6. (a) Define Quadratic residue codes of prime length and Extended Quadratic residue codes. Also illustrate each one by giving an example.8
 - (b) Prove that $F^{\perp} = \overline{N}$ and $N^{\perp} = \overline{F}$, where quadratic residue codes F and N of length $p \equiv 1 \pmod{4}$ over GF(s) generated by q(x) and n(x).
- 7. (a) Write normalized Hadamard matrix of order 8 and obtain the corresponding Hadamard codes.
 - (b) If M is a Hadamard matrix of order n and $T = \begin{pmatrix} M & M \\ M & -M \end{pmatrix}$, then prove that T is a Hadamard matrix of order 2n. Hence or otherwise write down a Hadamard matrix of order 4.

Unit IV

- 8. (a) If C is an [n, k, d] linear code over GF(q), then prove that C is MDS if and only if every n k columns of a parity check matrix H of C are linearly independent.8
 - (b) Prove that binary MDS codes are the trivial MDS codes.
- 9. (a) Let C be an [n, k] linear t-burst-error-correcting code. Then prove the following : 8
 - (i) no non-zero burst of length 2t or less can be codeword,
 - (ii) $n-k \geq 2t$.
 - (b) If C is an [n, k, d] MDS code, then show that any k symbols of the codewords may be taken as message symbols.

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